

QCLG Exercise sheet 8

Part 1. Quantum algorithm for the discrete logarithm problem

Let p be a prime number, $(\mathbb{Z}/p\mathbb{Z})^*$ the multiplicative group with elements in $\{1, \dots, p-1\}$. Let g be a generator of $(\mathbb{Z}/p\mathbb{Z})^*$. As a consequence, we have $g^{p-1} = 1 \pmod p$ and $\forall x \in (\mathbb{Z}/p\mathbb{Z})^*, \exists! y \in \{0, \dots, p-2\}, x = g^y \pmod p$.

The discrete log problem is the following: let $x \in (\mathbb{Z}/p\mathbb{Z})^*$, find $y \in \{0, \dots, p-2\}$ such that $x = g^y \pmod p$.

For any number $N \in \mathbb{N}$, we assume we have a perfect quantum unitary operation that performs the quantum Fourier transform QFT_N satisfying that $\forall k \in \{0, \dots, N-1\}$,

$$QFT_N(|k\rangle) = \frac{1}{\sqrt{N}} \sum_{a=0}^{N-1} \omega_N^{ak} |a\rangle.$$

where $\omega_N := e^{\frac{2i\pi}{N}}$.

Let V be the quantum unitary operator satisfying $V(|a\rangle |b\rangle |0\rangle) = |a\rangle |b\rangle |g^a x^{-b} \pmod p\rangle$. The goal of this exercise is to analyze the following quantum algorithm for the discrete logarithm problem.

Quantum algorithm for discrete logarithm.

We start with an input $x \in (\mathbb{Z}/p\mathbb{Z})^*$. We want to find $y \in \{0, \dots, p-2\}$ such that $x = g^y \pmod p$.

1. Start from three registers initialized at $|0\rangle$.
2. Apply QFT_{p-1} on each of the two first registers.
3. Apply the unitary operation V on all the registers.
4. Measure the third register in the computational basis *i.e.* the basis $\{|0\rangle, \dots, |p-1\rangle\}$. Let k be the output.
5. Apply the Fourier transform QFT_{p-1} on the first register and second register.
6. Measure both registers in the computational (standard) basis to get some outcome (l_1, l_2) .
7. Conclude

Let $|\phi_i\rangle$ the state of the algorithm after step i .

Exercise 1. Write $|\phi_1\rangle, |\phi_2\rangle, |\phi_3\rangle$. Let Z_c be the following set

$$Z_c := \{(a, b) \in \{0, \dots, p-2\} \times \{0, \dots, p-2\} : (a - by \pmod{p-1}) = c\}.$$

Show that $|\phi_3\rangle$ can be written

$$|\phi_3\rangle = \frac{1}{\sqrt{p-1}} \sum_{c \in \{0, \dots, p-2\}} \frac{1}{\sqrt{p-1}} \sum_{a, b \in Z_c} |a\rangle |b\rangle |g^c \pmod{p}\rangle.$$

Write therefore $|\phi_4\rangle$ as a function of c depending on the measured outcome g^c .

Exercise 2. For a fixed k (and hence c), write the state $|\phi_5\rangle$. Show that after step 6, we obtain a random couple (l_1, l_2) satisfying

$$l_1 y + l_2 = 0 \pmod{p-1}.$$

Conclude when l_1 is invertible in the ring $\mathbb{Z}/(p-1)\mathbb{Z}$. What would you do if this is not the case?

Part 2. Finding a hidden parabola in a function.

Suppose we are given a black-box function $f_{\alpha, \beta} : \mathbb{F}_p^2 \rightarrow \mathbb{F}_p$, where p is a prime, satisfying the promise that $f_{\alpha, \beta}(x, y) = f_{\alpha, \beta}(x', y')$ if and only if

$$\alpha x^2 + \beta x - y = \alpha x'^2 + \beta x' - y'.$$

for some unknown $\alpha \in \mathbb{F}_p$ and $\beta \in \mathbb{F}_p$. In other words, $f_{\alpha, \beta}$ is constant on the parabola

$$P_{\alpha, \beta, \gamma} := \{(x, y) \in \mathbb{F}_p^2 : y = \alpha x^2 + \beta x + \gamma\}$$

for any fixed $\gamma \in \mathbb{F}_p$, and distinct on parabolas corresponding to different values of γ . We have access to the unitary

$$O_{f_{\alpha, \beta}}(|x\rangle |y\rangle |z\rangle) = |x\rangle |y\rangle |z + f_{\alpha, \beta}(x, y)\rangle.$$

and our goal is to find α and β . Recall that for all $x \in \mathbb{F}_p$,

$$QFT_p(|x\rangle) = \frac{1}{\sqrt{p}} \sum_{y \in \mathbb{F}_p} \omega^{xy} |y\rangle.$$

where $\omega = e^{\frac{2i\pi}{p}}$. We consider the following procedure

Procedure 1

1. Start from three registers initialized at $|0\rangle$.
2. Apply QFT_p on each of the two first registers.
3. Apply the unitary operation $O_{f_{\alpha,\beta}}$ on all the registers.
4. Measure the third register in the computational basis *i.e.* the basis $\{|0\rangle, \dots, |p-1\rangle\}$.
5. Apply QFT_p on the second register and measure it.

Exercise 1. Show that after step 4 the above procedure creates the state

$$|\psi_4\rangle = \frac{1}{\sqrt{p}} \sum_{x \in \mathbb{F}_p} |x\rangle |\alpha x^2 + \beta x + \gamma\rangle.$$

for an unknown $\gamma \in \mathbb{F}_p$.

Exercise 2. Show that after step 5 the above procedure creates (up to a global phase) the state

$$|\psi_u\rangle = \frac{1}{\sqrt{p}} \sum_{x \in \mathbb{F}_p} \omega^{(\alpha x^2 + \beta x)u} |x\rangle.$$

for a known u . Why is u known?

Exercise 3. We apply the procedure twice and we construct the state $|\Phi\rangle = |\psi_u\rangle \otimes |\psi_{u'}\rangle$ for 2 known values $u, u' \in \mathbb{F}_p$. Write $|\Phi\rangle$. Show (using ancilla qubits), how to construct in time $O(\text{polylog}(p))$ the state

$$|\Omega_{u,u'}\rangle = \frac{1}{p} \sum_{x,x'} \omega^{\alpha(ux^2+u'x'^2)+\beta(ux+u'x')} |x\rangle |x'\rangle |ux^2 + u'x'^2\rangle |ux + u'x'\rangle$$

from $|\Phi\rangle$.

Exercise 4. We *assume* that, from $|\Omega_{u,u'}\rangle$, we know how to construct the state

$$|\xi\rangle = \frac{1}{p} \sum_{w_1, w_2 \in \mathbb{F}_p} \omega^{\alpha w_1 + \beta w_2} |w_1\rangle |w_2\rangle.$$

Think of a way to recover (α, β) from the state $|\xi\rangle$.